ThoughtWorks® Deutschland

Parallel Programming with Java 7

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Some theory
New Kids on the block
All hands on deck
Wrap up
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Some theory New Kids on the block All hands on deck Wrap vp 14:00 - 14:20 14:20 - 14:30 14:30 - 16:3016:30 - 16:45

Part 1 Some theory

Why parallel programming?

Grand challenge problems Modelling DNA structures Weather forecasting

Parallel programming splits problems into parts Parts are solved in parallel by multiple processors

No more free lunch



Microprocessor Transistor Counts 1971-2011 & Moore's Law

Source: http://de.wikipedia.org/wiki/Mooresches_Gesetz

Partitioning

Data partitioning is applied to data (aka domain decomposition) Functional decomposition concurrently executes independent functions Less common than data partitioning West Side Hwy Traffic Weather GW Bridge Traffic Canadian low pressure. Accident on I-95. Presidential motorcade, Hurricane in Gulf, Japanese butterfly. time of day. etc. etc..

7 Source: http://en.wikipedia.org/wiki/Functional_decomposition

Divide-and-conquer

Divide-and-conquer divides a problem into similar sub-problems Subproblems that have the same form as the original. Problem can be repeated recursively

Which way is faster?



Amdahl's law (1/3)

- **S** = speedup factor S = function(n,f) where
 - **n** = number of processors/cores
 - **f** = fraction of the computation that cannot be parallelised



Amdahl's law (2/3)

- S = speedup factor = function(n,f) where
 - n = number of processors/cores
 - f = fraction of the computation that cannot be parallelised

$$S(n) = \frac{\text{single processor execution time}}{n \text{ processor execution time}} = \frac{n}{1 + (n-1)f}$$
$$\lim_{n \to \infty} S(n) = \frac{1}{f}$$

Amdahl's law (3/3)



How to benefit from multiple CPUS/cores in Java

Implicitly - magic done under the hood by the JVM Explicitly - using Java 7's fork/join framework

This session focuses on the fork/join framework

Part 2 New Kids on the block

New classes on the block



Anatomy of the fork/join interaction (1/2)



Anatomy of the fork/join interaction (2/2)



Part 3 All hands on deck

How to implement a Fork Join Task generally

if (my portion of the work is small enough)
 do the work directly
else

split my work into two pieces and fork them
merge/join the results



Testing fork/join

- Assert that ...
- ... the complete task works as expected ... the core computation works as expected
- ... splitting the task does not corrupt input data
- ... merging results behaves as expected



A guided example: Calculating checksums

Setting up your environment

- ***** Sources and slides are on the stick or clone them from
 - https://github.com/fkoehler/TW-Fork-Join-Workshop
- ***** Follow the instructions in the README file
- ★ Get the JDK from
 - http://www.oracle.com/technetwork/java/javase/ downloads/jdk7-downloads-1637583.html



Exercise 1: Bucket sort

***** Implement a sequential and a parallel version of bucket sort

 Put elements into buckets that represent ranges



***** Sort and concatenate buckets

✗ Source: http://en.wikipedia.org/wiki/Bucket_sort



Exercise 1: Bucket sort as pseudo code

```
≮ function bucketSort(array, n) is
buckets ← new array of n empty lists
for i = 0 to (length(array)-1) do
insert array[i] into buckets[whichBucket(array[i])]
for i = 0 to n - 1 do
nextSort(buckets[i])
return the concatenation of buckets[0], ..., buckets[n-1]
```

✗ Source: http://en.wikipedia.org/wiki/Bucket_sort

Exercise 1: One way to realise bucket sort



Exercise 2: Tuning the parallel bucket sort

- ★ Buckets of equal width work well when the list to sort is evenly populated
- ★ When the distribution of elements is skewed, the parallel bucket sort may degrade (why?)
- ★ Run your parallel bucket sort with different kinds of lists (i.e. evenly populated ones and skewed ones)
- Improve the performance of your parallel bucket sort if required



Exercise 3: Merge sort

- Implement a sequential and a parallel version of merge sort
 - If the list is of length 0 or 1 the list is already sorted
 - Divide the unsorted list into two sublists of about half the size
 - Sort each sublist recursively by re-applying the merge sort
 - Merge the two sublists back into one sorted list



★ Source: http://en.wikipedia.org/wiki/Merge_sort

Exercise 3: Merge sort as pseudo code (1/2)

```
function merge sort(m)
    if length (m) \leq 1
        return m
    var list left, right, result
    var integer middle = length(m) / 2
    for each x in m up to middle
         add x to left
    for each x in m after or equal middle
         add x to right
    left = merge sort(left)
    right = merge sort(right)
    result = merge(left, right)
```

return result

Exercise 3: Merge sort as pseudo code (2/2)

```
function merge(left,right)
    var list result
    while length(left) > 0 or length(right) > 0
        if length(left) > 0 and length(right) > 0
            if first(left) ≤ first(right)
                append first(left) to result
                left = rest(left)
            else
                append first(right) to result
                right = rest(right)
        else if length(left) > 0
            append first (left) to result
            left = rest(left)
        else if length(right) > 0
            append first(right) to result
            right = rest(right)
    end while
    return result
```

29 Source: http://en.wikipedia.org/wiki/Merge_sort

Exercise 4: Searching strings in text files

- ***** Implement a parallel search for a string in a text file
- The search string can be assumed to contain no whitespace (i.e. no spaces or tabs)
- ***** Have a look at the following class:

com.thoughtworks.fjw.search.SimpleStringSearchTest

• The class contains examples for string based searching and reading from a text file

Exercise 5: Searching strings revisited

- ***** Revisit your solution for the previous exercise
- Drop the assumption so that search strings can contain whitespace
- Consider how dropping the assumption affects the way you partition the text to search



Part 4 Wrap up

B. Wilkinson & M. Allen: Parallel Programming, Prentice Hall, New Jersey 1999



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Backup slides

Amdahl's law (3/3)



Amdahl's law (3/3)



Thread state diagram

